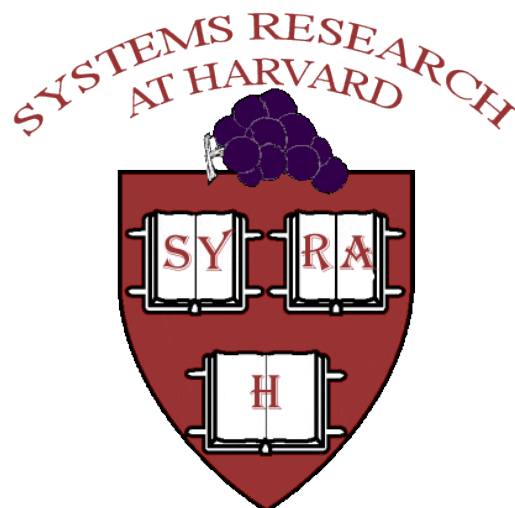


# Web Scale Data Management: An Historical Perspective



**Margo Seltzer**

March 1, 2018

# Outline

- In the beginning ...
- The heyday of RDBMS
- The rebirth of key/value stores
- Key/Value stores today: NoSQL
- NoSQL & key/value use cases

# In the Beginning

- Where beginning equals 1960's
- Computers
  - Centralized systems
  - Spiffy new data channels let CPU and IO overlap.
  - Persistent storage is on drums.
  - Buffering and interrupt handling done in the OS.
  - Making these systems fast is becoming a research focus.
- Data
  - What did data look like?

# Organizing Data: ISAM

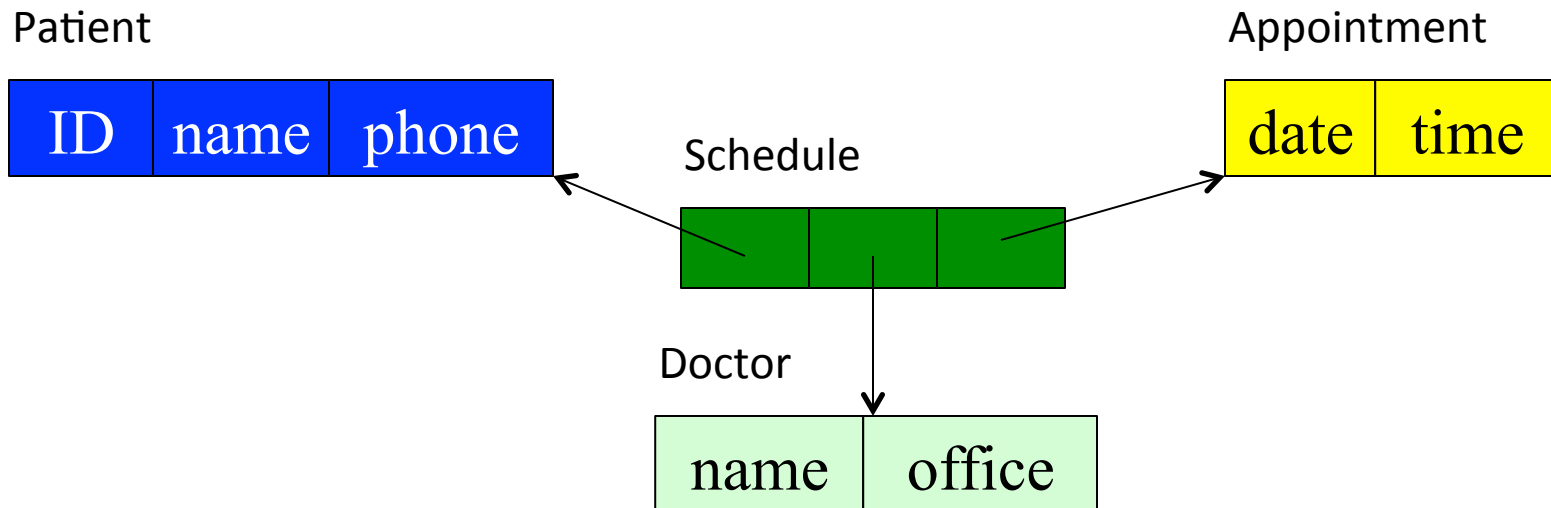
- Indexed sequential access method
  - Pioneered by IBM for its mainframes
  - Fixed length records
  - Each record lives at a specific location
  - Rapid record access even on a sequential medium (tape)
- All indexes are ‘secondary’
  - Allow key lookup, but ...
  - Do not correspond to physical organization
  - Key is to build small, efficient index structures
- Fundamental access method in COBOL

# Organizing Data: The Network Model

- Early data management systems represented data using something called a network model.
- Data are represented by collections of records (*today we would call those key/data pairs*).
- Relationships among records are expressed via links between records (today we can think of those links as pointers).
- Applications interacted with data by navigating through it:
  - Find a record
  - Follow links
  - Find other records
  - Repeat

# The Network Model: Inside Records

- Records composed of attributes (think fields)
- Attributes are single-valued.
- Links connect exactly two records.
- Represent N-way relationships via link records



# The Relational Model: The Competition

- The Network Model had some problems
  - Applications had to know the structure of the data
  - Changing the representation required a massive rewrite
  - Fundamentally: the physical arrangement was tightly coupled to the application and the application logic.
- 1968: Ted Codd proposes the relational model
  - Decouple physical representation from logical representation
  - Store “records” as “tables”
  - Replace links with implicit joins among tables
- The big question: could it perform?

# Outline

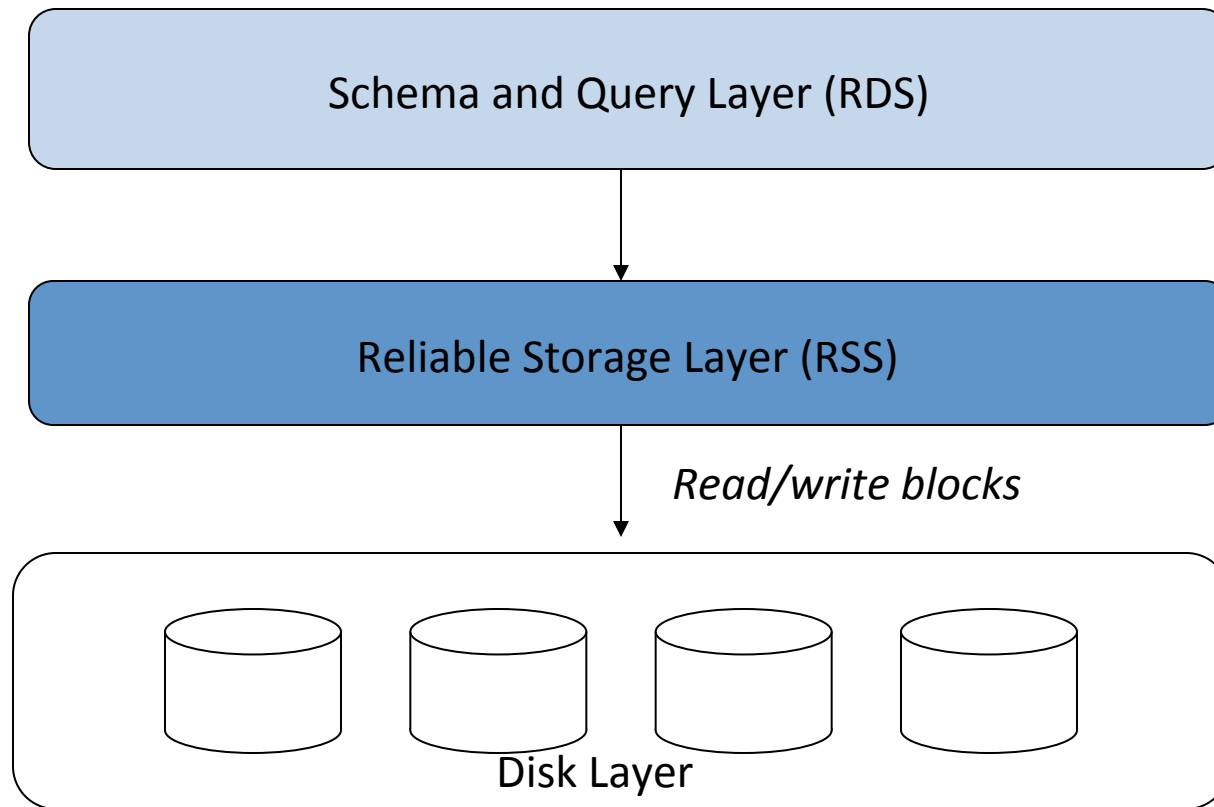
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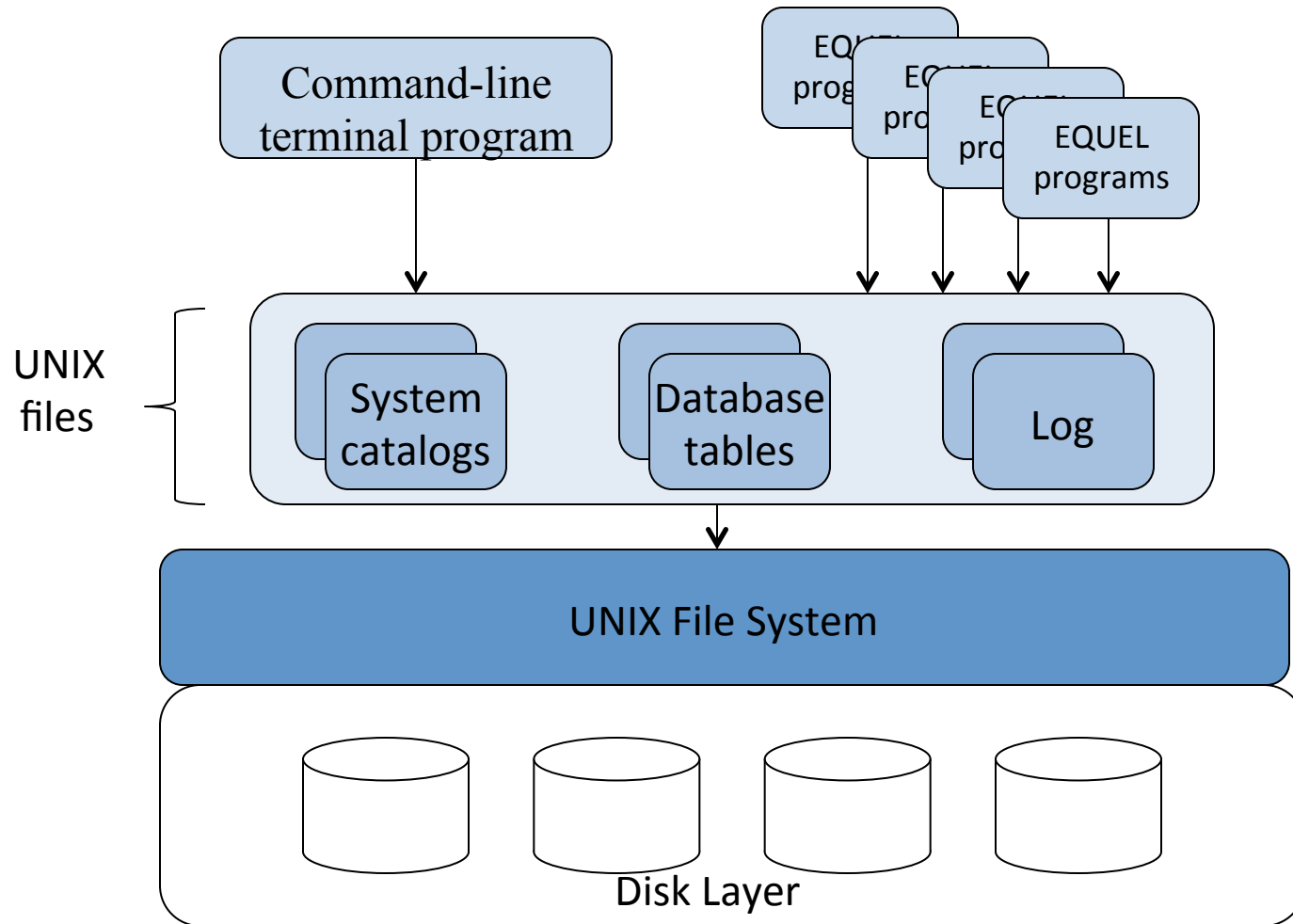
# The heyday of RDBMS

- After Codd's paper, much debate ensued, but two groups set out to turn an idea into software.
  - IBM: System/R
  - U.C.Berkeley: Ingres
- Both were research projects to explore the feasibility of implementing the relational model.
- Both were hugely successful and had enormous impact.
- However, at their core, you'll notice something interesting ...

# The Design of System R



# The Design of Ingres



# The Key/Value Store Within

- All these relational systems and most data management systems have hidden deep inside some sort of key/value storage engine.
- For years, the “conventional wisdom” was to hide those KV stores deep underneath a query language and schema level.
- Major exception was COBOL, which continued to use ISAM.

# RDBMS matured

- Lots of new features (SQL-86, SQL-89, SQL-92, SQL:1999, SQL:2003, SQL:2008, SQL:2100?)
  - Triggers
  - Stored procedures
  - Report generators
  - Rules
  - ...
- Incorporated new data models:
  - Object-relational systems
  - XML
- Became enormously large, complex systems requiring expert administration.

# RDBMS: The Good and the Bad

- RDBMS Advantages:
  - Declarative query language that decoupled physical and logical layout.
  - (Relative) ease of schema modification.
  - Great support for ad hoc queries.
- RDBMS Disadvantages:
  - Pay for the overhead of query processing even if you don't have complex queries.
  - Require DBA for tuning and maintenance.
  - Nearly always pay IPC to access database server.
  - Require schema definition.
  - Not great at managing hierarchical or other complex relationships.

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# The Age of the Internet

- New classes of applications emerged
  - Search
  - Authentication (LDAP)
  - Email
  - Browsing
  - Instant messaging
  - Web servers
  - Online retail
  - Public key management



# These Applications were Different

- They weren't supporting ad hoc query mechanisms
  - Query set specified by the application
  - Users interacted with the applications not the data
- They had fairly simple schemas
- No need for fancy reports
- Performance was critical

*All the bells and whistles of the big relational players weren't delivering necessary functionality, but were costing overhead.*

# The Emergence of Key/Value Stores

- In 1997, Sleepycat Software, introduced the first commercial key value store (now Oracle Berkeley DB).
- Berkeley DB was/is:
  - Embedded: links directly in an application's address space
  - Fast: avoids query parsing and optimization
  - Scalable: runs from handheld to data center
  - Flexible: trade off functionality and performance
  - Reliable: recovers from application and system failure

# Key/Value versus RDBMS

	Key-value Databases	Standard RDBMS
Pros	<ul style="list-style-type: none"><li>• Sometimes the “good enough” solution for simple data schemas</li><li>• Much smaller footprint</li><li>• Shorter execution path</li><li>• More efficient, fewer OS resources</li><li>• Usually no client/server</li><li>• No application data mapping</li></ul>	<ul style="list-style-type: none"><li>• Rich SQL Support</li><li>• Data typing</li><li>• Data relationship management</li><li>• Dynamic data relationships</li><li>• Procedural languages (stored procedures)</li><li>• Parallel query execution</li><li>• Security, Encryption</li><li>• Centralized management</li></ul>
Cons	<ul style="list-style-type: none"><li>• Application-centric schema control</li><li>• High Availability often missing (not BDB)</li><li>• Higher level abstraction APIs often missing (not BDB)</li><li>• Limited RDBMS integration</li></ul>	<ul style="list-style-type: none"><li>• Much larger engine, higher overhead, client/server overhead, general purpose RDBMS solution</li><li>• Higher administrative burden</li><li>• Higher TCO</li></ul>

# From local to distributed

- As the web matured, data volume and velocity and customer demand grew exponentially.
- Exceeded single system capacity.
- Enter a new era of scalability demands.
- Infrastructure migrates from large-scale SMP to clusters of commodity blades.

# Outline

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- **Key/Value stores today: NoSQL**
  - Framing
  - Technology
- NoSQL & key/value use cases

# Enter NoSQL

- Web service providers (e.g., Google, Amazon, Ebay, Facebook, Yahoo) began pushing existing data management solutions to their limits.
- Introduced sharding: Split the data up across multiple hosts to spread out load.
- Introduced replication:
  - Allow access from multiple sites
  - Increase availability
- Relax consistency: Didn't always need transactional semantics

# What is NoSQL?

- Not-only-SQL (2009)
- While RDBMs have typically focused on transactions and the ACID properties, NoSQL focuses on BASE:
  - **B**asic **A**vailability: Use **replication** to reduce the likelihood of data unavailability and **sharding** to make any remaining failures partial
  - **S**oft state: Allow data to be inconsistent and relegate designing around such inconsistencies to application developers.
  - **E**ventually consistent: Ensure only that at some future point in time the data assumes a consistent state

# Why NoSQL?

- Three common problems
  1. Unprecedented transaction volumes
  2. Critical need for low latency access
  3. 100% Availability
- Hardware shift: From massive SMP to blades
- Changing how we deal with the CAP theorem
  - The CAP theorem states that you can't have all three of:
    - **C**onsistency: All nodes see the same data at the same time.
    - **A**vailability: Every request receives a success/fail response
    - **P**artition tolerance: The system continues to operate despite arbitrary message loss.
  - RDBMs (transactional systems) typically choose CA
  - NoSQL typically chooses AP or CP



# NoSQL Evolution

- 1997: Berkeley DB released providing transactional key/value store.
- 2001: Berkeley DB introduces replication for high availability.
- 2006: Google publishes Chubby and BigTable papers.
  - Each system provides scalable, data management for loosely coupled systems
- 2007: Amazon publishes Dynamo paper
  - DHT-based eventually consistent key/value store
- 2008+: Multiple OSS projects to produce BigTable/Dynamo knockoffs
  - HBase: Apache Hadoop project implementation of BigTable (2008)
  - CouchDB: Erlang-based document store, MVCC and versioning (2008)
  - Cassandra: Write-optimized, column-oriented, secondary indices (2009)
  - MongoDB: JSON-based document store, indexing, sharded (2009)
- 2009+: Commercialization
  - Companies emerge to support open source products (e.g., DataStax/Cassandra, Basho/Riak, Cloudera/HBase, 10Gen/MongoDB )
  - Companies develop commercial offerings (Oracle, Citrusleaf)

# NoSQL Technology

- Some form of distribution and replication:
  - Distributed hash tables (DHTs)
  - Key partitioning
  - Replication in underlying storage or in NoSQL
- Relational or key/value store for underlying storage engine.
- Second generation systems building custom write-optimized engines.
  - Log-structured merge trees (LSM)

# NoSQL Design Space

- Local node storage system
- Distribution
- Data Model
- Consistency Model
  - Eventual consistency
  - No consistency
  - Transactional consistency

# Local Storage Systems

- Native KVstore
  - Oracle NoSQL Database: Berkeley DB Java Edition
  - Basho: Riak's Bitcask, a log-structured hash table
  - Amazon: Dynamo, BDB Data Store, BDB JE (or MySQL)
- Log-Structured Merge Trees
  - LevelDB
- Custom
  - BigTable (and clones): Tablet servers on SSTables exploiting GFS-like systems.

# Distribution

- Three main questions:
  - How do you partition data
  - How many copies do you keep
  - Are all copies equal
- How do you partition data
  - Key partitioning
  - Hash partitioning (often called sharding)
  - Geographic partitioning (also called sharding)

# More Distribution

- How many copies and how?
  - Use the underlying file system (GFS, HDFS)
  - Three is good; five is better; for hot things sometimes go to many.
- Copy Equality:
  - Single-Master
    - MongoDB, Oracle NoSQL Database,
  - Multi-Master/Masterless
    - Riak, CouchDB, Couchbase

# NoSQL Data Models

- Common to most offerings:
  - Denormalization: some data values are duplicated to provide superior query performance.
  - No joins
- Key/Value: Little or no schema, minimal range query support
  - Oracle NoSQL DB, DynamoDB, Couchbase, Riak
- BigTable Column Family
  - Hbase, Cassandra
- Document-stores
  - MongoDB, CouchDB

# Data Model: Key/Value

- Keys and data are both opaque byte strings.
- Designed for point queries
  - Typically no notion of a range query
  - If iteration exists, it's usually not in key-order
- Examples:
  - DynamoDB
  - LevelDB
  - Riak
  - Tokyo/Kyoto Cabinet
  - Oracle NoSQL Database
- Extensions:
  - Oracle NoSQL: Major/Minor keys, batching
  - LevelDB: Atomic batches



# Data Model: Column Family (1)

- Columns are grouped into *column families*.
- Column families:
  - Are typically stored together
  - Can have different columns for each row
  - Can have duplicate items in any column
- No schema or type enforcement
  - All data are treated as byte strings
- Indexed by rows
  - Rows are grouped into tablets
  - No secondary indexes

# Data Model: Column Family (2)

Column Families

Timestamps

Keys	Name		Address					
1	First Margo	Last Seltzer	No 3	Street Millstone Lane	City Lincoln	State MA	Zip 01773	2000
			No 394	Street East Riding Dr	City Carlisle	State MA	Zip 01741	1993
			PO 65		City Sonyea	State NY	Zip 14556	1961
3	Title Lady	Last Gaga	City Hollywood			State CA	Zip 90027	2008
4	Last Madonna		New York		Los Angeles		London	2000

versions

# Data Model: Document

- Key/Value store where the value is a document with structure.
- Document store understands the structure of documents:
  - JSON, XML, BSON, PDF, Doc
- Sometimes documents can have sub-documents.
- Key attribute of a document store is that it lets you search within documents as well as searching for documents.

# NoSQL Consistency

- The CAP theorem (Eric Brewer): it is impossible for a distributed computer system to simultaneously provide:
  - Consistency: All nodes see the same data at the same time
  - Availability: Every request receives a response about whether it succeeded
  - Partition tolerance: The system continues to operate despite arbitrary message loss or failure of parts of the system
- Originally posed as a conjecture by Brewer, later precisely proven by Gilbert and Lynch.
- It's an interesting history to read!
- Implications for NoSQL: Pick 2
  - Some systems pick CP
  - Other systems pick AP
  - Few systems pick CA
- Why?

# Forms of Consistency

- Consistent/Partitionable Systems
  - BigTable, Hbase, HypterTable, MongoDB, Redis, MemcacheDB
- Available/Partitionable Systems
  - Cassandra, SimpleDB, CouchDB, Riak, TokyoCabinet, Dynamo
- Eventual Consistency
  - There exist multiple definitions. The most popular one is due to Vogels: The storage system guarantees that if no new updates are made to the object, eventually (after the inconsistency window closes) all accesses will return the last updated value.
  - You can be AP and eventually consistent.
- Variable: through careful configuration choose how consistent you want to be:
  - Cassandra: read/write to one, quorum, all
  - Oracle NoSQL DB: Can be CP when setting ack policy for simple majority, else is AP.
  - Amazon Dynamo: Picking read/write quorums trades off performance and degree of inconsistency

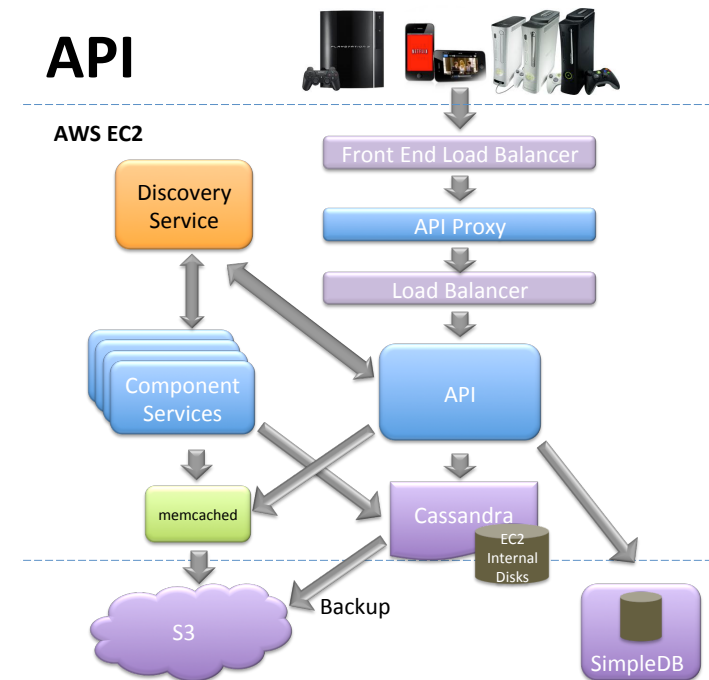
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# Cassandra Use Case: Netflix

Adrian Cockcroft <http://www.hpts.ws/sessions/GlobalNetflixHPTS.pdf>

- In 2011 Netflix moved from centralized data centers to an out-sourced, distributed cloud-based system.
- Replaced centralized database with Cassandra for:
  - Customers
  - Movies
  - History
  - Configuration
- Why?
  - No need for schema changes
  - High performance
  - Scalability



# HBase Use Case: Facebook Messages

Kannan Muthukkaruppan:

<http://www.hpts.ws/sessions/StorageInfraBehindMessages.pdf>

- Facebook integrated messaging, chat, email and SMS under a single new message framework.
- Why?
  - High write throughput
  - Good read performance
  - Horizontal scalability
  - Strong consistency
- What is in Hbase?
  - Small messages
  - Message metadata
  - Search index



# MongoDB Use Case: Viper Media

<http://nosql.mypopescu.com/post/16058009985/mongodb-at-viber-media-the-platform-enabling-free>

- Viper uses MongoDB in a central service to which mobile clients connect to route messages to other clients.
- Why?
  - Scalability
  - Redundancy
- What is in MongoDB?
  - Variable length documents
  - Dictionary indexes

# Wrapping Up

- When should you be considering NoSQL?
  - Scalability
  - Low latency
  - Redundancy
  - No adhoc queries
  - Joins easily implementable in the application
  - Straight forward key structure
- A couple of sites I found pretty interesting:
  - List of NoSQL Databases: <http://nosql-database.org>
  - NoSQL Data Modeling: <http://highlyscalable.wordpress.com/2012/03/01/nosql-data-modeling-techniques/>
  - Recent Performance Comparison: <http://www.networkworld.com/cgi-bin/mailto/x.cgi?pagetosend=/news/tech/2012/102212-nosql-263595.html>