

Imperative, OO and Functional Languages

- A "C" program is ...
 - a web of assignment statements, interconnected by control constructs which describe the time sequence in which they are to be executed.
- In Java programming,
 - "objects" are sent "messages".
- In "pure" LISP there is only ...
 - the evaluation of an expression by function application
 - instead of "executing a program," LISP evaluates a symbolic expression (s-expr)

Why Learn LISP ???

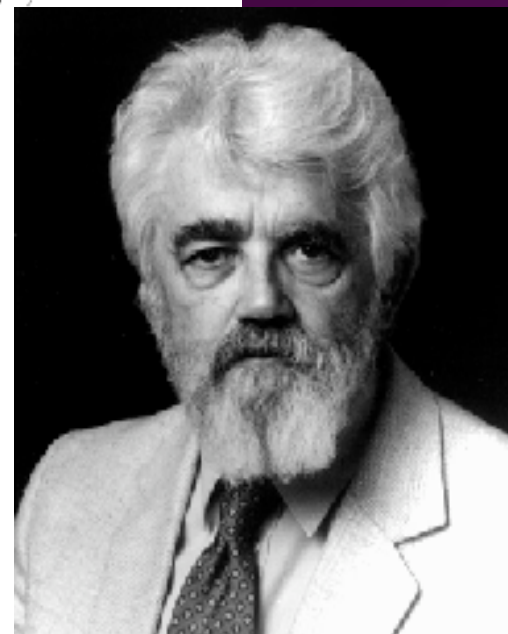
- "Lisp is worth learning for the profound enlightenment experience you will have when you finally get it; that experience will make you a better programmer for the rest of your days, even if you never actually use Lisp itself a lot."
- Eric Raymond, "How to Become a Hacker"
- For examples of some companies that use LISP, see <http://www.paulgraham.com/apps.html>

Where Does LISP Come From?

Recursive Functions of Symbolic Expressions and Their Computation by Machine, Part I

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April 1960



1 Introduction

A programming system called LISP (for LIST Processor) has been developed for the IBM 704 computer by the Artificial Intelligence group at M.I.T. The system was designed to facilitate experiments with a proposed system called the Advice Taker, whereby a machine could be instructed to handle declarative as well as imperative sentences and could exhibit “common sense” in carrying out its instructions. The original proposal [1] for the Advice Taker was made in November 1958. The main requirement was a programming system for manipulating expressions representing formalized declarative and imperative sentences so that the Advice Taker system could make deductions.

Introduction to LISP

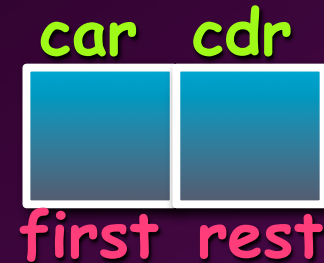
- LISP syntax: symbolic expressions
 - atoms
 - lists of atoms
 - lists of s-exprs
- Grammar for symbolic expressions:
 - $\langle \text{s-expr} \rangle ::= \langle \text{atom} \rangle \mid \langle \text{list} \rangle$
 - $\langle \text{list} \rangle ::= (\langle \text{s-expr} \rangle^*)$
- (name-or-description-of-a-function
arg₁ arg₂ ... arg_n)

Evaluation of S-EXPRs

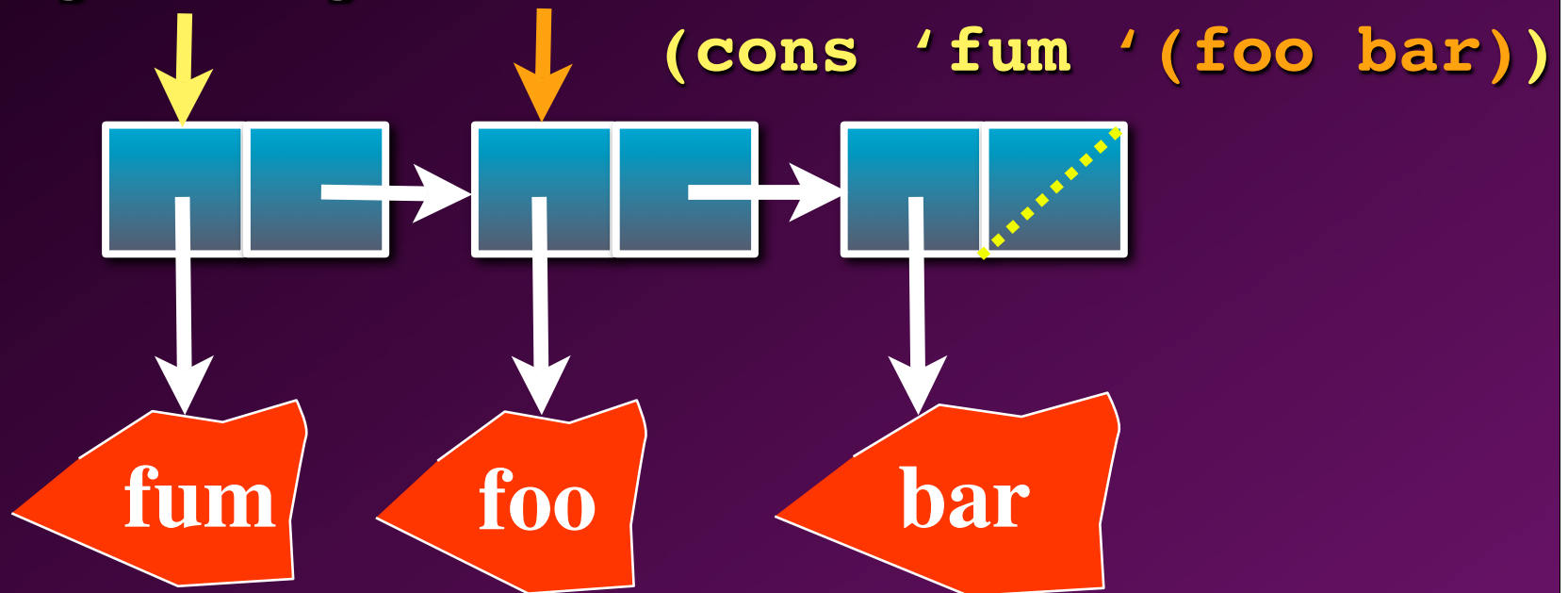
- Parentheses must be taken seriously!
- Quoting inhibits evaluation (EVAL does the opposite!)
- **NIL** is both a list and an atom
- Defining your own functions using DEFUN
 - area of a circle
- Other useful functions: **IF**, **+**, *****, **=**
- Defining a recursive function: **factorial**

Internal Representation

- LISP lists are stored as linked-lists of records with CAR and CDR fields



- Diagramming list structure:



Other LISP Functions

- Assignment (side-effects) using **SETF**, **SET**
- Lisp Manipulation Functions
 - **CAR, CDR, CONS, LIST, APPEND**
- Predicates
 - **EQ, EQL, EQUAL, ATOM, LISTP, CONSP, NULL, ZEROP, PLUSP, MINUSP, EVENP, ODDP, NUMBERP, SYMBOLP, BOUNDP, >, <, <=, >=, =, /=**
 - Define a function to recursively compute the length of a list

Predicates

- Summary of commonly confused primitive predicates:

| | 'A | '(A) | nil |
|---------|-----|------|-----|
| atom | t | nil | t |
| lisp | nil | t | t |
| consp | nil | t | nil |
| null | nil | nil | t |
| symbolp | t | nil | t |

Defining 2 More Functions

- Swap the first and second elements of a list.

➤ e.g., `(swap '(A B C D))`  `(B A C D)`

➤

```
(defun swap (lst)
  (cons '(cadr lst)
        (cons (car lst) (cddr lst))))
```

➤ Note: this builds a new list with some structure shared with L

- Compute the "next even integer" following n

➤

```
(defun next-even (n)
  (if (evenp n) n (1+ n)))
```

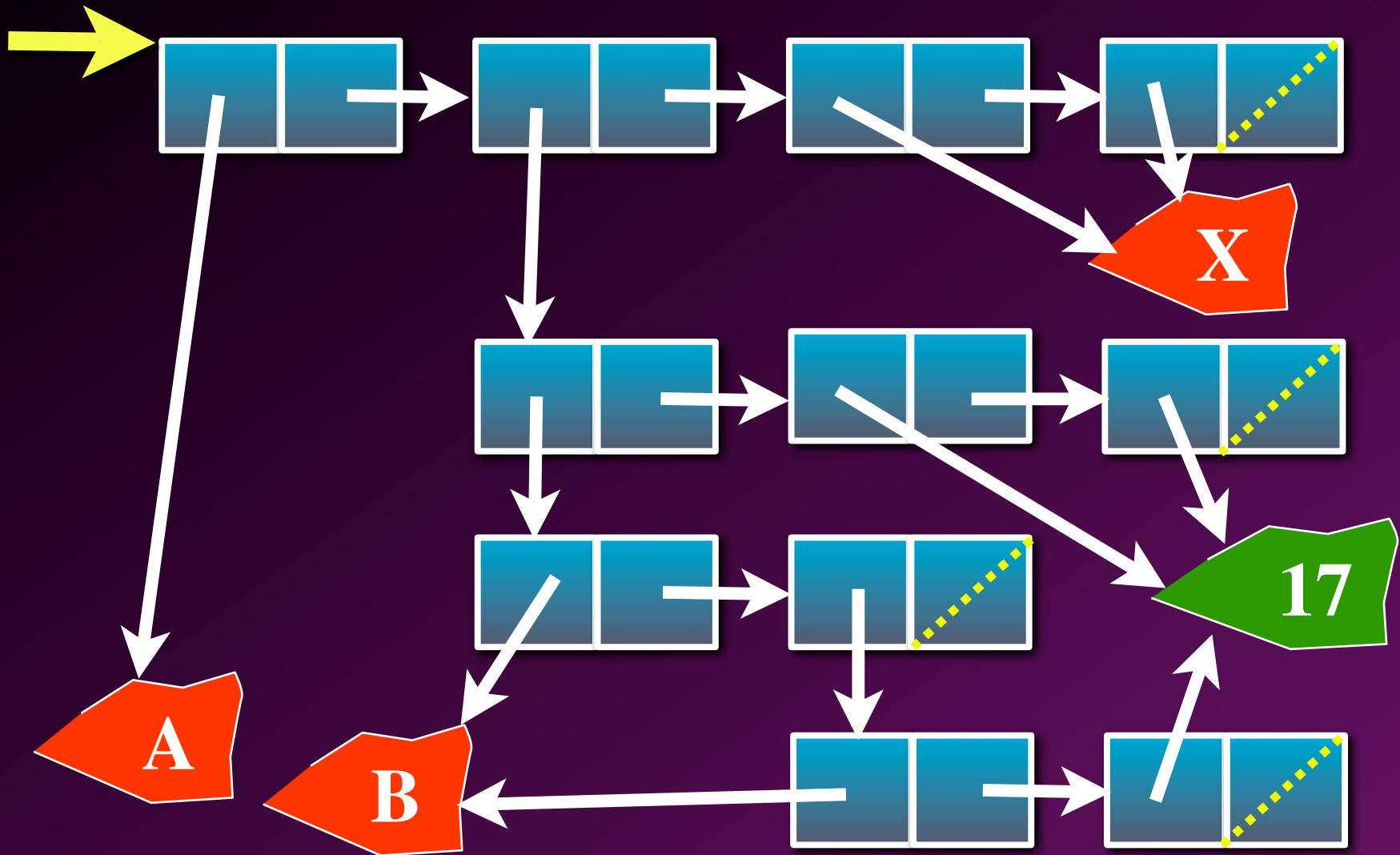
The Most General Conditional

$(cond \quad (test_1 \quad s\text{-}expr_{11} \quad s\text{-}expr_{12} \quad \dots \quad s\text{-}expr_{1n_1})$
 $(test_2 \quad s\text{-}expr_{21} \quad s\text{-}expr_{22} \quad \dots \quad s\text{-}expr_{2n_2})$
 \dots
 \dots
 $(test_k \quad s\text{-}expr_{k1} \quad s\text{-}expr_{k2} \quad \dots \quad s\text{-}expr_{kn_k}))$

Recursion on List Structures

- **CAR** and **CDR** take lists apart
- the analog of “subtract 1” (when doing induction on a number) is taking the **CDR** of a list
- Our version of the built-in MEMBER function:
 - ```
(defun memb (element lis)
 (cond ((null lis) nil)
 ((eq (car lis) element) lis)
 (t (memb element (cdr lis)))))
```
- Recursion in “two directions” — function OCCURS
  - ```
(defun occurs (element lis)
  (cond ...
```

Example of Nested List



Nameless Functions via LAMBDA

- LISP programs are conceived and written with a mathematical rigor, based upon the formalisms of "recursive function theory" and the "lambda calculus."
 - Consider $y + (x * y)$ for the values 3 and 4
 - Clarify using the "Lambda notation" of Alonzo Church.
 - In LISP, we use a similar notation
- (MAPCAR F L)
 - F is a function of one parameter
 - L is a list $(x_1 \ x_2 \ \dots \ x_n)$
 - Produces a list $(y_1 \ y_2 \ \dots \ y_n)$ where $y_i = (F \ x_i)$

Procedural Abstraction

- **SUM-INTEGERS** computes

$$\sum_{n = \text{first}}^{\text{last}} n$$

- **SUM-SQUARES** computes

$$\sum_{n = \text{first}}^{\text{last}} n^2$$

- Now make the function-of-n itself a third parameter:

- **SUM-TERMS** computes

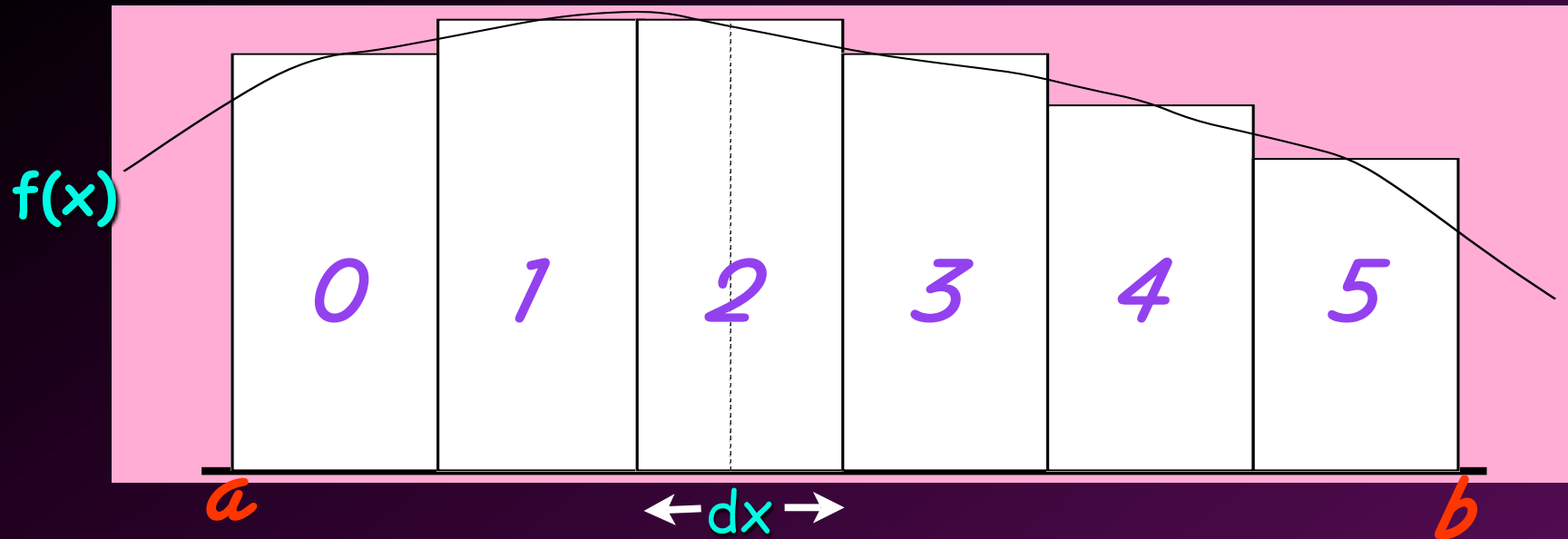
$$\sum_{n = \text{first}}^{\text{last}} \text{term-fn}(n)$$

- Consider now the infinite series

$$\frac{\pi}{8} = \frac{1}{1 * 3} + \frac{1}{5 * 7} + \frac{1}{9 * 11} + \dots$$

and define **PI-TERM(n)** to compute the above

Integration by Summation



$$\blacksquare \sum_{n=0}^{(b-a)/dx - 1} f(a + n \cdot dx + dx/2) \cdot dx$$

$$\blacksquare dx \cdot \sum_{n=0}^{(b-a)/dx - 1} f(a + n \cdot dx + dx/2)$$

Symbolic Pattern Matching

- Another kind of search problem: in a linear list of words (symbols, whatever) to discover specified patterns.
 - Although LISP itself has no built-in pattern-matching, it's a good implementation language for such a function.
 - `(match pattern data)` will return T or NIL
 - The `pattern` may contain "wildcard" variables such as `?` (stand for one symbol) and `*` (stands for a sequence of 0 or more symbols)

Pattern Matching, part 2

■ "Wildcard" examples

➤ **A ? B** matches A A B
but not A B
nor A B C

➤ **A * B** matches A A B
and A B
and A X Y Z B

➤ *** X * Y** matches any sequence containing
both X and Y in that order

Additional "Wildcards"

- **?variable** matches a single atom, and assigns that atom to **variable**
- ***variable** matches a sequence of ≥ 0 atoms, and assigns a list of that sequence to **variable**
- Example from Doctor program:
 - `(cond ((match '(I am worried *blah-blah) userInput)
 (princ (append '(How long have you been worried)
 blah-blah)))`